

## CLAIMS

Having thus described our invention, what we claim as new, and desire to secure by Letters Patent is:

- 1 Sub 1. A method for mapping a valid stack up to a destination program counter,  
2 comprising:  
3 A6 mapping a path of control flow on the stack from any start point in a  
4 selected method to the destination program counter; and  
5 simulating stack actions for executing bytecodes along said path.

- 1 2. The method of claim 1 wherein the step of mapping a path of control flow on the  
2 stack comprises:  
3 processing a first linear bytecode sequence until the control flow is interrupted;  
4 and recording unprocessed targets from any branches in the first linear bytecode  
5 sequence for future processing.

- 1 3. The method of claim 2 wherein the step of mapping a path of control flow on the  
2 stack further comprises:  
3 processing an additional bytecode linear sequence until the control flow is  
4 interrupted; and  
5 recording unprocessed targets from any branches in the additional linear bytecode  
6 sequence for future processing, where the destination program counter was not  
7 reached during an earlier processing of a linear bytecode sequence.

- 1 4. The method of claim 2 wherein the step of processing any linear bytecode  
2 sequence comprises:

3 determining if a bytecode in said any linear bytecode sequence is a breakpoint  
4 with a pointer to bytecode data; and  
5 replacing the breakpoint with the bytecode data.

1 5. The method of claim 3 wherein the step of processing any linear bytecode  
2 sequence comprises:

3 determining if a bytecode in said any linear bytecode sequence is a breakpoint  
4 with a pointer to bytecode data; and  
5 replacing the breakpoint with the bytecode data.

6. The method of claim 1 wherein the step of simulating stack actions executing the bytecodes along the path further comprises generating a virtual stack.

1            7. The method of claim 6, further comprising:

2        encoding the virtual stack as a bitstring and storing the bitstring at a selected  
3        destination for use in memory management operations.

1 8. The method of claim 7, wherein the step of storing the bitstring comprises storing  
2 the bitstring to the selected method as compiled on a heap.

1        9. The method of claim 7, wherein the step of storing the bitstring comprises storing  
2        the bitstring to a pre-allocated area on the stack.

1        10. The method of claim 1 wherein the step of simulating stack actions executing the  
2        bytecodes along the path further comprises:

3            inserting pre-determined stack actions for bytecodes maintaining the control flow  
4            in the selected method; and

030509Z  
010300Z

Sub 63  
calculating stack actions for bytecodes transferring the control flow from the selected method.

Sub A7  
1 11. A method for mapping a Java bytecode stack up to a destination program counter  
2 comprising:  
3 mapping a path of control flow on the stack from any start point in a selected  
4 method to the destination program counter; and  
5 simulating stack actions for executing bytecodes along said path.

1 12. The method of claim 11 wherein the step of mapping a path of control flow on  
2 the stack comprises:  
3 processing a first linear bytecode sequence until the control flow is interrupted;  
4 and recording unprocessed targets from any branches in the first linear bytecode  
5 sequence for future processing.

Sub 63  
09324005-061099  
1 13. The method of claim 12 wherein the step of mapping a path of control flow on  
2 the stack further comprises:  
3 processing an additional bytecode linear sequence until the control flow is  
4 interrupted; and  
5 recording unprocessed targets from any branches in the additional linear bytecode  
6 sequence for future processing, where the destination program counter was not  
7 reached during an earlier processing of a linear bytecode sequence.

1 14. The method of claim 12 wherein the step of processing any linear bytecode  
2 sequence comprises:  
3 determining if a bytecode in said any linear bytecode sequence is a breakpoint  
4 with a pointer to bytecode data; and

5 replacing the breakpoint with the bytecode data.

1 15. The method of claim 13 wherein the step of processing any linear bytecode  
2 sequence comprises:

3 determining if a bytecode in said any linear bytecode sequence is a breakpoint  
4 with a pointer to bytecode data; and  
5 replacing the breakpoint with the bytecode data.

Sub C3 1 16. The method of claim 11 wherein the step of simulating stack actions executing the  
2 bytecodes along the path further comprises generating a virtual stack.

09329558-061099 1 17. The method of claim 16 further comprising:  
2 encoding the virtual stack as a bitstring and storing the bitstring at a selected  
3 destination for use in memory management operations.

1 18. The method of claim 17, wherein the step of storing the bitstring comprises  
2 storing the bitstring to the selected method as compiled on a heap.

1 19. The method of claim 17, wherein the step of storing the bitstring comprises  
2 storing the bitstring to a pre-allocated area on the stack.

1 20. The method of claim 11 wherein the step of simulating stack actions executing the  
2 bytecodes along the path further comprises:

3 inserting pre-determined stack actions for bytecodes maintaining the control flow  
4 in the selected method; and

5 calculating stack actions for bytecodes transferring the control flow from the  
6 selected method.

1 21. A computer-readable memory for storing the instructions for use in the execution  
in a computer of the method of claim 1.

1 22. A computer readable memory for storing the instructions for use in the execution  
2 in a computer of the method of claim 11.

23. A program storage device readable by a machine, tangibly embodying a program  
of instructions executable by the machine to perform method steps for mapping a  
valid stack up to a destination program counter, said method steps comprising:

mapping a path of control flow on the stack from any start point in a selected  
method to the destination program counter; and

simulating stack actions for executing bytecodes along said path,

wherein the step of mapping a path of control flow on the stack comprises:

processing a first linear bytecode sequence until the control flow is interrupted;  
and

recording unprocessed targets from any branches in the first linear bytecode  
sequence for future processing, and

where the destination program counter was not reached during an earlier  
processing of a linear bytecode sequence,

processing an additional bytecode linear sequence until the control flow is  
interrupted; and

recording unprocessed targets from any branches in the additional linear bytecode  
sequence for future processing.